

Classical gates

State of classical computers: sequence of 0 and 1.

Transformations are called gates.

Def.: a classical (logical) gate is $g: \{0,1\}^n \rightarrow \{0,1\}$ (elementary), $g: \{0,1\}^m \rightarrow \{0,1\}^m$ (extended).

Ex.: NOT:

Notation: \oplus is $+ \pmod{2}$

$$\text{NOT}(x_1) = 1 \oplus x_1;$$

$$\text{AND}: (x_1, x_2) \mapsto x_1 \cdot x_2;$$

$$\text{XOR}: (x_1, x_2) \mapsto x_1 \oplus x_2;$$

$$\text{OR}: (x_1, x_2) \mapsto x_1 \oplus x_2 + x_1 \cdot x_2;$$

$$\text{TOFFOLI}: \{0,1\}^3 \rightarrow \{0,1\}^3.$$

$$(x_1, x_2, x_3) \mapsto (x_1, x_2, x_1 \cdot x_2 \oplus x_3)$$

Def.: a gate is reversible if its associated function is a bijection.

$$\text{Remark: } \text{TOF}(x_1, x_2, 0) = (x_1, x_2, \text{AND}(x_1, x_2));$$

$$\text{TOF}(1, x_2, x_3) = (1, x_2, \text{XOR}(x_2, x_3));$$

$$\text{TOF}(1, 1, x_3) = (1, 1, \text{NOT}(x_3)).$$

More useful gates: ID: $x_i \mapsto x_i$;

$$\text{FALSE}: x_1 \mapsto 0;$$

$$\text{TRUE}: x_1 \mapsto 1;$$

$$\text{COPY}: x_1 \mapsto (x_1, x_1).$$

(FANOUT)

Rules for combining gates

Let g_1, \dots, g_k gates. Define $F(g_1, \dots, g_k)$ the set of gates that can be constructed from g_1, \dots, g_k according to the following rules:

$$(1) g_1, \dots, g_k \in F(g_1, \dots, g_k);$$

$$(2) \text{padding}, P_{j_1, \dots, j_k; j_1, \dots, j_k}^{(m)}: \{0,1\}^m \rightarrow \{0,1\}^{m+k};$$

$$(x_1, \dots, x_m) \mapsto (x_1, \dots, x_{j_1-1}, y_1, x_{j_1}, x_{j_1+1}, \dots)$$

$$(3) \text{restrictions and reordering}: r_{j_1, \dots, j_k}^{(m)}: \{0,1\}^m \rightarrow \{0,1\}^k, k \leq m, \text{ when it makes sense}$$

$$(x_1, \dots, x_m) \mapsto (x_{j_1}, \dots, x_{j_k})$$

$$(4) \text{composition of gates}: h_1, h_2 \in F(g_1, \dots, g_k) \Rightarrow h_1 \circ h_2 \in F(g_1, \dots, g_k);$$

$$(5) \text{cartesian products}: \text{"", "} \in \text{""} \Rightarrow h_1 \times h_2 \in \text{""}.$$

Def.: a set of gates g_1, \dots, g_k is universal if \forall gate $g, g \in F(g_1, \dots, g_k)$.

Theorem: the Toffoli gate is universal and reversible. Proof: book. \square

Quantum Gates

$$\rightarrow \mathbb{H} \cong \mathbb{C}^2$$

Def.: a quantum n -gate is a unitary operator $U: \mathbb{H}^{\otimes n} \rightarrow \mathbb{H}^{\otimes n}$.

The qubits on which we apply quantum gates form a quantum register.

Ex.: • identity $\text{II} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$;

• phase factor $M(\alpha) = e^{i\alpha} \text{II} = \begin{pmatrix} e^{i\alpha} & 0 \\ 0 & e^{i\alpha} \end{pmatrix}$;

• phase shift $P(\alpha) = |0\rangle \langle 0| + e^{i\alpha} |1\rangle \langle 1| = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\alpha} \end{pmatrix}$;

• Q NOT (Pauli X) $\text{Q}_X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$;

• Pauli Y $\sigma_Y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}$;

• Pauli Z $\sigma_Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$;

• Hadamard $H = \frac{\text{Q}_X + \text{Q}_Z}{\sqrt{2}} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$;

• spin rotation $D_m(\alpha) = \dots$.

Measurement of observable A: $\text{A} \rightarrow \lambda$ ($\lambda = \text{measured value}$).

• C NOT (controlled quantum NOT) $\Lambda^1(X) = |0\rangle \langle 0| \otimes \text{II} + |1\rangle \langle 1| \otimes X$

(variants: $\Lambda_1(X) = \text{II} \otimes |0\rangle \langle 0| + X \otimes |1\rangle \langle 1|$,

$\Lambda^{10}(X) = |0\rangle \langle 0| \otimes X + |1\rangle \langle 1| \otimes \text{II}$);

• in general: if V is a quantum (unary) gate, controlled V is

$\Lambda^1(V) = |0\rangle \langle 0| \otimes \text{II} + |1\rangle \langle 1| \otimes V = \begin{pmatrix} \text{II} & 0 & 0 & 0 \\ 0 & \text{II} & 0 & 0 \\ 0 & 0 & \text{II} & 0 \\ 0 & 0 & 0 & \text{II} \end{pmatrix}$;

• swap gate $\text{S} = |00\rangle \langle 00| + |11\rangle \langle 11| + |11\rangle \langle 11| + |01\rangle \langle 10| + |10\rangle \langle 01|$,

can be written as $\text{S} = \text{II} \otimes \text{II} + \text{II} \otimes \text{II} + \text{II} \otimes \text{II} + \text{II} \otimes \text{II} + \text{II} \otimes \text{II}$.

Rules for combining quantum gates

Given $U_1, \dots, U_k, U_j \in \mathcal{U}(\mathbb{H}^{\otimes n_j})$, the set $F(U_1, \dots, U_k)$ of gates that can be obtained from U_1, \dots, U_k is defined by:

$$(1) U_1, \dots, U_k \in F(U_1, \dots, U_k);$$

$$(2) \text{II}^{\otimes m} \in F(U_1, \dots, U_k);$$

$$(3) V_1, V_2 \in F(U_1, \dots, U_k) \Rightarrow V_1 V_2 \in F(U_1, \dots, U_k);$$

$$(4) \text{if } V_1 \in \mathcal{U}(\mathbb{H}^{\otimes n_1}), V_2 \in \mathcal{U}(\mathbb{H}^{\otimes n_2}), V_1, V_2 \in F(U_1, \dots, U_k) \Rightarrow V_1 \otimes V_2 \in F(U_1, \dots, U_k).$$

$\{U_1, \dots, U_k\}$ is a universal set of gates if any unitary operator on $\mathbb{H}^{\otimes m}$ belongs to $F(U_1, \dots, U_k)$.

Theorem: $\{M, D_X, D_Z, \text{C NOT}\}$ is a universal set.