

$\varphi: A_* \rightarrow B_*$ morfismo di complessi di g.a.:
 $\forall i \quad A_i \xrightarrow{\varphi_i} B_i \quad d'_i \circ \varphi_i = \varphi_{i+1} \circ d_i$.

$\begin{array}{ccc} |d_i| & \rightarrow & |d'_i| \\ A_{i+1} \xrightarrow{\varphi_{i+1}} B_{i+1} & & \end{array}$

$\varphi_i(\ker d_i) \subset \ker d'_i, \quad \varphi_i(\operatorname{Im} d_{i-1}) \subset \operatorname{Im} d'_{i-1} \Rightarrow$
 $\Rightarrow \varphi \text{ induce } \varphi_*: H^*(A_*) \rightarrow H^*(B_*) \text{ omomorfismo di g.a.}$

$0 \rightarrow A_* \xrightarrow{\varphi} B_* \xrightarrow{\psi} C_* \rightarrow 0$ esatta di complessi di g.a. \Rightarrow
 $\Rightarrow 0 \rightarrow H^0(A_*) \xrightarrow{\varphi_0} H^0(B_*) \xrightarrow{\psi_0} H^0(C_*) \xrightarrow{\delta_0} H^1(A_*) \xrightarrow{\varphi_1} H^1(B_*) \rightarrow$
 $\xrightarrow{\psi_1} H^1(C_*) \xrightarrow{\delta_1} H^2(A_*) \rightarrow \dots$ esatta di g.a..

$\begin{array}{ccccccc} 0 & \rightarrow & A_i & \xrightarrow{\varphi_i} & B_i & \xrightarrow{\psi_i} & C_i & \rightarrow & 0 \\ & & \downarrow d_i & & \downarrow d'_i & & \downarrow d''_i & & \\ 0 & \rightarrow & A_{i+1} & \xrightarrow{\varphi_{i+1}} & B_{i+1} & \xrightarrow{\psi_{i+1}} & C_{i+1} & \rightarrow & 0 \end{array} \quad \begin{array}{l} \text{diagram} \\ \text{chasing} \end{array}$

Vogliamo $\varphi_i: H^i(C_*) \rightarrow H^{i+1}(A_*)$, si fa come a ETA.

Sia $0 \rightarrow L \xrightarrow{\Phi} F \xrightarrow{\Psi} G \rightarrow 0$ esatta di prefasci di g.a. su X ,
il ricoprimento aperto di X , $\forall p$
 $0 \rightarrow \check{C}^p(U, L) \xrightarrow{\Phi_p} \check{C}^p(U, F) \xrightarrow{\Psi_p} \check{C}^p(U, G) \rightarrow 0$ esatta,

dove dato $\sigma \in \check{C}^p(U, L)$, $\sigma = (\ell_{\alpha_0, \dots, \alpha_p})$,

$\Phi_p(\sigma) = (\phi_{U_{\alpha_0, \dots, \alpha_p}}(\ell_{\alpha_0, \dots, \alpha_p})), \quad \ell_{(U_{\alpha_0, \dots, \alpha_p})} \quad \Psi_p \text{ analogo.}$

Troviamo la succ. esatta lunga per i $\check{H}^*(U, \cdot)$.

Poiché \lim_{\leftarrow} è esatto, abbiamo

$0 \rightarrow \check{H}^0(X, L) \rightarrow \check{H}^0(X, F) \rightarrow \check{H}^0(X, G) \xrightarrow{\delta_0} \check{H}^1(X, L) \rightarrow \check{H}^1(X, F) \rightarrow \check{H}^1(X, G) \rightarrow \dots$ esatta.

Se $0 \rightarrow L \xrightarrow{\varphi} \mathcal{F} \xrightarrow{\psi} \mathcal{G} \rightarrow 0$ è esatta di fasci di g.a. su X ,
abbiamo $0 \rightarrow \Gamma(X) \rightarrow \Gamma(\mathcal{F}) \rightarrow \Gamma(\mathcal{G}) \xrightarrow{\psi(\Gamma(L))} 0$ esatta di prefasci di g.a.
su $X \Rightarrow \dots \rightarrow \check{H}^i(X, \mathcal{L}) \rightarrow \check{H}^i(X, \mathcal{F}) \rightarrow \check{H}^i(X, \mathcal{G}) \rightarrow \dots$ esatta lunga.

Se X è paracompatto, $\check{H}^i(X, G) = \check{H}^i(X, \text{Sheaf}(G))$,

$\text{Sheaf}(G) = \text{Coh}^0(G) \cong G$.

Def.: un fascio \mathcal{F} su X si dice FIACCO se $\forall \emptyset \neq U \subset X$ aperto la
restrizione $\rho_U^X: \Gamma(X, \mathcal{F}) \rightarrow \Gamma(U, \mathcal{F})$ è suri. (es.: i $\text{Can}^i(\mathcal{F})$
sono fiacchi). Ciò implica: $\forall \emptyset \neq V \subset U \subset X$ aperti, ρ_V^U è suri..

Prop.: $0 \rightarrow \mathcal{L} \xrightarrow{\varphi} \mathcal{F} \xrightarrow{\psi} \mathcal{G} \rightarrow 0$ esatta di fasci di g.a. su X , \mathcal{L} fiacco;

- 1) $0 \rightarrow \Gamma(\mathcal{L}) \xrightarrow{\varphi} \Gamma(\mathcal{F}) \xrightarrow{\psi} \Gamma(\mathcal{G}) \rightarrow 0$ esatta di prefasci di g.a. su X ;
- 2) \mathcal{F} fiacco $\Leftrightarrow \mathcal{G}$ fiacco.

Dim.: 1) Vogliamo: $\forall \emptyset \neq U \subset X$ aperto $0 \rightarrow \Gamma(U, \mathcal{L}) \xrightarrow{\varphi_U} \Gamma(U, \mathcal{F}) \xrightarrow{\psi_U} \Gamma(U, \mathcal{G}) \rightarrow 0$
esatta, ci basta ψ_U suri.. $\lambda \in \Gamma(U, \mathcal{G})$, sia T l'insieme delle
coppie (U', t') dove $\emptyset \neq U' \subset U$ aperto e $t' \in \Gamma(U', \mathcal{F})$
t.c. $\psi_U t' = \psi_{U'}(t') = \lambda|_{U'}$. λ si solleva localmente $\Rightarrow T \neq \emptyset$.

Semiordiniamo T : $(U', t') \leq (U'', t'')$ se $U' \subset U''$ e t'' estende
 t' ($t''|_{U'} = t'$). Allora $\exists (V, t) \in T$ massimale.

Vogliamo $V = U$, per assurdo sia $x \in U \setminus V$. $\exists x \in U_x \subset U$ aperto e
 $t_x \in \Gamma(U_x, \mathcal{F})$ t.c. $\psi_{U_x}(t_x) = \lambda|_{U_x}$. Su $V \cap U_x$,

$\psi_{V \cap U_x}(t|_{V \cap U_x}) = \lambda|_{V \cap U_x} = \psi_{U_x}(t_x) = (\psi_{U_x}(t_x)|_{V \cap U_x})$, cioè
 $\psi(t) - \psi(t_x) = 0 \Rightarrow$ a meno di restrinmare U_x , su $V \cap U_x$

$t - t_x$ proviene da una sezione di \mathcal{L} $\ell \in \Gamma(U_x \cap V, \mathcal{L})$
 $((t - t_x)|_{V \cap U_x} = \psi_{U_x \cap V}(\ell))$. Estendiamo ℓ a una sezione

ℓ' di \mathcal{L} su U_x $\ell' \in \Gamma(U_x, \mathcal{L})$, $\ell'|_{U_x \cap V} = \ell \Rightarrow$
 $\Rightarrow t_x + \psi_{U_x}(\ell') \in \Gamma(U_x, \mathcal{F})$ coincide con t su $U_x \cap V$ e
si reincolla su t dando una sezione t'' di \mathcal{F} su $V \cup U_x$

che solleva λ , assurdo. Allora $V = U \Rightarrow \psi_U(t) = \lambda$.

2) $0 \rightarrow \Gamma(X, \mathcal{L}) \rightarrow \Gamma(X, \mathcal{F}) \rightarrow \Gamma(X, \mathcal{G}) \rightarrow 0$

$0 \rightarrow \Gamma(U, \mathcal{L}) \rightarrow \Gamma(U, \mathcal{F}) \rightarrow \Gamma(U, \mathcal{G}) \rightarrow 0$

commutativo con righe e prima colonna esatta.

$\Rightarrow 0 \rightarrow \dots \rightarrow \dots \rightarrow 0 \quad \text{diagram chasing.}$

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